

Distant decimals of π

Yves Bertot, Laurence Rideau, Laurent Théry

June 2018

- ▶ Arithmetic-geometric means: an algorithm with $3n$ divisions
- ▶ Arithmetic-geometric means: an algorithm with only one division
- ▶ A few facts about the correctness of individual digits
- ▶ The BBP formula

The Arithmetic-Geometric computation

- ▶ A pair of two sequences of real numbers
- ▶ $a_0 = 1$ $b_0 = x$ $a_{n+1} = \frac{a_n + b_n}{2}$ $b_{n+1} = \sqrt{a_n b_n}$
- ▶ $\int_{-\infty}^{\infty} \frac{dt}{\sqrt{(1^2+t^2)(x^2+t^2)}} = \int_{-\infty}^{\infty} \frac{dt}{\sqrt{(a_n^2+t^2)(b_n^2+t^2)}}$
- ▶ a_n and b_n converge very fast to the same value (a_n is larger)
- ▶ $a_n - b_n < A^{-2^n}$
- ▶ Number of known digits doubles at every step

AGM for π

- ▶ Define $f(x)$ as the arithmetic-geometric mean of 1 and x
- ▶ Key property $\pi = 2\sqrt{2} \frac{f(1/\sqrt{2})^3}{f'(1/\sqrt{2})}$
- ▶ Mathematical proofs based on elliptic integrals
 - ▶ Improper integrals (Coquelicot)
 - ▶ Needed extensions (Chasles with improper integrals)
 - ▶ Lots of formulas with squares and square roots slightly beyond Presburger automated proofs

Borwein&Borwein algorithm

► Define $y_n(x) = \frac{a_n(x)}{b_n(x)}$ $z_n(x) = \frac{b'_n(x)}{a'_n(x)}$

► Easy consequence $1 + y_n = 2 \frac{a_{n+1} b_{n+1}^2}{a_n b_n^2}$ $1 + z_n = 2 \frac{a'_{n+1}}{a'_n}$

$$y_0(x) = \frac{1}{x} \quad y_{n+1} = \frac{1 + y_n}{2\sqrt{y_n}} \quad z_1 = \frac{1}{\sqrt{x}} \quad z_{n+1} = \frac{1 + z_n y_n}{(1 + z_n)\sqrt{y_n}}$$

► $\pi_n = (2 + \sqrt{2}) \prod_{i=1}^n \frac{1 + y_i}{1 + z_i}$ at $x = \frac{1}{\sqrt{2}}$

convergence rate

$$0 \leq \pi_{n+1} - \pi \leq 8\sqrt{2} \times 531^{-2^n}$$

Brent and Salamin algorithm

- ▶ Define $c_n = \frac{1}{2}(a_{n-1} - b_{n-1})$
- ▶ $\frac{b'_{n+1}}{b_{n+1}} - \frac{b'_n}{b_n} = \frac{b_n}{2a_n} \left(\frac{a_n}{b_n} \right)'$ $\frac{b'_{n+1}}{b_{n+1}} = \frac{b'_1}{b_1} - \sqrt{2} \sum_{k=1}^{n-1} 2^k c_k^2$
- ▶ $\pi'_n = \frac{4a_n^2}{1 - \sum_{k=1}^{n-1} 2^{k-1} (a_{k-1} - b_{k-1})^2}$

convergence rate

- ▶ First approximation $|\pi'_{n+1} - \pi| \leq 68 \times 531^{-2^{n-1}}$
- ▶ Coarser than Borwein&Borwein: $\pi_{n+1} - \pi \leq 8\sqrt{2} \times 531^{-2^n}$
- ▶ Improvement by studying $|\pi'_{n+2} - \pi'_{n+1}|$

$$|\pi'_{n+1} - \pi| \leq (132 + 384 \times 2^n) \times 531^{-2^n}$$

- ▶ For computing 10^6 decimals of π both π'_{19} and π_{19} are enough
 - ▶ $132 + 384 \times 2^{19} \leq 2^{28}$
- ▶ Each algorithm computes n square roots, but π_n computes $3n$ divisions, π'_n only half-sums and one full division.

Rounding errors

- ▶ Approximating real computations using fixed-point computations (rounding towards 0)
- ▶ Take the same program code, replace operations
- ▶ How do rounding errors propagate?
- ▶ Amazingly $|\boxed{y_n} - y_n| \leq 2ulp$ and $|\boxed{z_n} - z_n| \leq 4ulp$
- ▶ $|\boxed{\pi_n} - \pi_n| \leq (21 * n + 2)ulp$
- ▶ $|\boxed{\pi'_n} - \pi'_n| \leq (160(\frac{3}{2})^{n+1} + 80 * 3^{n+1} + 100)ulp$

Total interval for Brent-Salamin

- ▶ In the end $10^{10^6+k} \times \pi'_{19}$ is within 2^{40} of $10^{10^6+k} \times \pi$
- ▶ It remains to choose a suitable value of k
- ▶ It is more efficient to compute $2^{\lfloor \log_2(10) \rfloor \times 10^6} \times \pi'_{19}$
- ▶ Paper available at <https://hal.inria.fr/hal-01582524>
- ▶ Code and instructions <https://www-sop.inria.fr/marelle/distant-decimals-pi/>
- ▶ Includes a C implementation of the Borwein algorithm, on top of MPFR.

The BBP formula

Work done in Coq by Laurence Rideau and Laurent Théry

$$\pi = \sum_{i=0}^{\infty} \frac{1}{16^i} \left(\frac{4}{8i+1} - \frac{2}{8i+4} - \frac{1}{8i+5} - \frac{1}{8i+6} \right).$$

- ▶ Mathematical justification using Riemann integrals
- ▶ commuting integration and infinite sum: thanks to Coquelicot

Computation of a single (hexa-decimal) digit

- ▶ Computing the d^{th} hexadecimal digit of π
- ▶ $\lfloor 16^{d-1}\pi \rfloor (\text{mod } 16)$
- ▶ Choose a precision p
- ▶ Compute separately each of the sums, taking the modulo right-away
- ▶ Example: $\sum_{i=0}^{\infty} \lfloor \frac{2^p 16^{d-1-i} \times 4}{8i+1} \rfloor (\text{mod } 2^p)$

Intuitive example

- ▶ Transposing in base 10, what is the second digit (on the right of the dot) of

$$\sum_{i=0}^{\infty} \frac{1}{10^i} \frac{1}{2i+1}$$

- ▶ $1 + 0.03333\dots + 0.00200\dots + 0.00014\dots + 0.00001\dots + \dots$

Computation of a single digit (continued)

- ▶ Only a finite number of terms in the sum, approximately $d + p/4$ terms
 - ▶ No modulo is needed for the last $p/4$ terms
- ▶ Use integer division: uncertainty bounded by 1
- ▶ Accumulated uncertainty is $d + p/4 + 1$
- ▶ We need the accumulated uncertainty to be small wrt. 2^{p-4}
- ▶ Final result is partial: when the value modulo 2^p is closer to 2^p than the accumulated uncertainty the digit is not known.